## Pencil-based algorithms for the tensor rank decomposition are not stable

Paul Breiding (TU Berlin, MPI MiS)
Carlos Beltrán (Universidad de Cantabria)
Nick Vannieuwenhoven (KU Lueven)

personal-homepages.mis.mpg.de/breiding/
juliahomotopycontinuation.org

Hitchcock (1927) introduced the tensor rank decomposition or CPD for tensors $\mathcal{A} \in \mathbb{R}^{n_{1} \times \cdots \times n_{d}}$ :

$$
\mathcal{A}=\sum_{i=1}^{r} \mathbf{a}_{i}^{1} \otimes \mathbf{a}_{i}^{2} \cdots \otimes \mathbf{a}_{i}^{d}
$$



The rank of a tensor is the minimum number of rank-1 tensors of which it is a linear combination.

Notation
$\mathcal{S}:=\{\mathcal{A} \mid \operatorname{rank}(\mathcal{A})=1\}$ are the rank-one tensors.
$\sigma_{r}:=\{\mathcal{A} \mid \operatorname{rank}(\mathcal{A}) \leq r\}$ are the tensors of rank at most $r$.

## A direct algorithm for order-3 tensors

State-of-the art algorithms compute the CPD by applying optimization algorithms on the goal function $\frac{1}{2}\left\|\mathcal{A}-\sum_{i=1}^{r} \mathcal{A}_{i}\right\|^{2}$.

However, in some cases, the CPD of third-order tensors can be computed directly via a generalized eigendecomposition.

For simplicity, assume that $\mathcal{A} \in \mathbb{R}^{n \times n \times n}$ is of rank $n$. Say

$$
\mathcal{A}=\sum_{i=1}^{n} \mathbf{a}_{i} \otimes \mathbf{b}_{i} \otimes \mathbf{c}_{i} .
$$

The steps are as follows.

1. Choose a matrix $Q \in \mathbb{R}^{n \times 2}$ with orthonormal columns $\mathbf{q}_{1}, \mathbf{q}_{2}$. 2. Compute the multilinear multiplication

2. The two 3 -slices $X_{1}$ and $X_{2}$ of $x$ are

$$
X_{j}=\sum_{i=1}^{n}\left\langle\mathbf{q}_{j}, \mathbf{c}_{i}\right\rangle \mathbf{a}_{i} \otimes \mathbf{b}_{i}=A \operatorname{diag}\left(\mathbf{q}_{j}^{T} C\right) B^{T}
$$

where $A=\left[\mathbf{a}_{i}\right] \in \mathbb{R}^{n \times n}$ and likewise for $B$ and $C$. Hence, $X_{1} X_{2}^{-1}$ has the following eigenvalue decomposition:

$$
X_{1} X_{2}^{-1}=A \operatorname{diag}\left(\mathbf{q}_{1}^{T} C\right) \operatorname{diag}\left(\mathbf{q}_{2}^{T} C\right)^{-1} A^{-1}
$$

from which $A$ can be found as the matrix of eigenvectors.
4. By a 1-flattening

we find

$$
\mathcal{A}_{(1)}:=\sum_{i=1}^{n} \mathbf{a}_{i}\left(\mathbf{b}_{i} \otimes \mathbf{c}_{i}\right)^{T}=A(B \odot C)^{T}
$$

where $B \odot C:=\left[\mathbf{b}_{i} \otimes \mathbf{c}_{i}\right]_{i} \in \mathbb{R}^{n^{2} \times n}$. Computing

$$
A \odot\left(A^{-1} \mathcal{A}_{(1)}\right)^{T}=A \odot(B \odot C)=\left[\mathbf{a}_{i} \otimes \mathbf{b}_{i} \otimes \mathbf{c}_{i}\right]_{i}
$$

solves the tensor decomposition problem.

Let's perform an experiment in Tensorlab v3.0 with this decomposition algorithm.

Create a rank- 25 random tensor of size $25 \times 25 \times 25$ :
\% Ut\{i\} has as columns the i-th factors
>> Ut\{1\} = randn $(25,25)$;
>> Ut\{2\} = randn $(25,25)$;
>> Ut\{3\} = randn $(25,25)$;
\% generate the full tensor
>> A = cpdgen(Ut);

Compute $\mathscr{A}$ 's decomposition and its distance to the input decomposition, relative to the machine precision $\epsilon \approx 2 \cdot 10^{-16}$ :

```
>> Ur = cpd_gevd(A, 25);
>> E = kr(Ut) - kr(Ur);
>> norm( E(:), 2 ) / eps
ans =
    8.6249e+04
```

What happened?

Let us look more closely at the computational problem:

- The input is a tensor $\mathcal{A} \in \sigma_{25} \subset \mathbb{R}^{25 \times 25 \times 25}$ of rank 25 .
- The output is the tuple $\left(\mathbf{a}_{i} \otimes \mathbf{b}_{i} \otimes \mathbf{c}_{i}\right)_{i=1}^{25} \in \mathcal{S}^{\times 25}$.

Let $f: \sigma_{25} \rightarrow \mathcal{S}^{\times 25}$ be the function that maps a tensor to its decomposition. Then, what we observed was

$$
\frac{\left\|f(\mathscr{A})-f\left(\mathfrak{A}^{\prime}\right)\right\|}{\left\|\mathcal{A}-\mathcal{A}^{\prime}\right\|} \approx 8 \cdot 10^{4}
$$

with $\left\|\mathcal{A}-\mathfrak{A}^{\prime}\right\| \approx 2 \cdot 10^{-16}$.

The condition number quantifies the worst-case sensitivity of $f$ to perturbations of the input.


When $f$ is differentiable: $\kappa[f](x)=\left\|\mathrm{d}_{x} f\right\|$.
Problem: There is no function $f: \sigma_{25} \rightarrow \mathcal{S}^{\times 25}$ that maps a tensors to its decomposition.

## A short detour through algebraic geometry

If the set of rank-1 tensors $\left\{\mathcal{A}_{1}, \ldots, \mathscr{A}_{r}\right\}$ is uniquely determined given the rank-r tensor $\mathcal{A}=\mathcal{A}_{1}+\cdots+\mathcal{A}_{r}$, then we call $\mathcal{A}$ an $r$-identifiable tensor.

Note that matrices are never $r$-identifiable, because

$$
M=\sum_{i=1}^{r} \mathbf{a}_{i} \otimes \mathbf{b}_{i}=A B^{T}=\left(A X^{-1}\right)\left(B X^{T}\right)^{T}
$$

for every invertible $X$. In general, these factorizations are different.

Kruskal (1977) gave a famous sufficient condition for proving the $r$-identifiability of third-order tensors.

More recently $r$-identifiability was studied in algebraic geometry. This is a natural framework because the set of rank-1 tensors

$$
\mathcal{S}^{\mathbb{C}}:=\left\{\mathbf{a}^{1} \otimes \mathbf{a}^{2} \otimes \cdots \otimes \mathbf{a}^{d} \mid \mathbf{a}^{k} \in \mathbb{C}^{n_{k}} \backslash\{0\}\right\}
$$

is the smooth projective Segre variety.
The set of tensors of rank bounded by $r$,

$$
\sigma_{r}^{\mathbb{C}}:=\left\{\mathcal{A}_{1}+\cdots+\mathcal{A}_{r} \mid \mathcal{A}_{i} \in \mathcal{S}\right\},
$$

is the Zariski-open constructible part of the projective $r$-secant variety of the Segre variety.

In the words of algebraic geometry:
$\sigma_{r}^{\mathbb{C}}$ is generically $r$-identifiable, if the addition map:

$$
\begin{aligned}
\Phi_{r}: \mathcal{S}^{\mathbb{C}} \times \cdots \times \mathcal{S}^{\mathbb{C}} & \rightarrow \sigma_{r}^{\mathbb{C}} \\
\left(\mathcal{A}_{1}, \ldots, \mathcal{A}_{r}\right) & \mapsto \mathcal{A}_{1}+\cdots+\mathcal{A}_{r}
\end{aligned}
$$

is of degree $r$ !.
Let $n_{1} \geq \cdots \geq n_{d}$ and

$$
r_{\mathrm{cr}}=\frac{n_{1} \cdots n_{d}}{1+\sum_{i=1}^{d}\left(n_{i}-1\right)}, \quad r_{\mathrm{ub}}=n_{2} \cdots n_{d}-\sum_{k=2}^{d}\left(n_{k}-1\right) .
$$

Conjectured general rule:

$$
\begin{array}{ll}
r \geq r_{\mathrm{cr}} \text { or } d=2 & \rightarrow \\
\text { not gen. } r \text {-identifial } \\
n_{1}>r_{\mathrm{ub}} \text { and } r \geq r_{\mathrm{ub}} & \rightarrow \\
\text { not gen. } r \text {-identifial } \\
\text { none of foregoing and } r<r_{\mathrm{cr}} & \rightarrow \\
\text { gen. } r \text {-identifiable; }
\end{array}
$$

see Chiantini, Ottaviani, Vannieuwenhoven (2014) for a proof in the case $n_{1} \cdots n_{d} \leq 15000$.

The real case is more involved because now

$$
\sigma_{r}:=\sigma_{r}^{\mathbb{R}}:=\left\{\mathcal{A}_{1}+\cdots+\mathcal{A}_{r} \mid \mathcal{A}_{i} \in \mathcal{S}^{\mathbb{R}}\right\}
$$

is a semi-algebraic set.

Qi, Comon, and Lim (2016) showed that if $\sigma_{r}^{\mathbb{C}}$ is generically $r$-identifiable, then it follows that the set of real rank- $r$ tensors with multiple complex CPDs is contained in a proper Zariski-closed subset of $\sigma_{r}^{\mathbb{R}}$.

In this sense, $\sigma_{r}^{\mathbb{R}}$ is also generically $r$-identifiable.

In the following we abbreviate $\mathcal{S}:=\mathcal{S}^{\mathbb{R}}, \sigma_{r}:=\sigma_{r}^{\mathbb{R}}$ and we will assume that $\sigma_{r}^{\mathbb{C}}$ is generically $r$-identifiable.

## Defining the tensor decomposition map

To define the tensor decomposition map, we analyze the real addition map:

$$
\begin{aligned}
\Phi_{r}: \mathcal{S} \times \cdots \times \mathcal{S} & \rightarrow \mathbb{R}^{n_{1} \times \cdots \times n_{d}} \\
\left(\mathcal{A}_{1}, \ldots, \mathscr{A}_{r}\right) & \mapsto \mathcal{A}_{1}+\cdots+\mathscr{A}_{r}
\end{aligned}
$$

Note that the domain and codomain are smooth manifolds.
The idea is to define $\Phi_{r}$ on the quotient of $\mathcal{S}^{\times r}$ by the symmetric group and restricting the domain of $\Phi_{r}$ to a Zariski-open smooth submanifold. Then, $\Phi_{r}$ restricts to a diffeomorphism onto its image.

Let $\mathbf{n}=\left(n_{1}, \ldots, n_{d}\right)$.
Let $\mathcal{M}_{r ; \mathbf{n}} \subset \mathcal{S}^{\times r}$ be the set of tuples of $n_{1} \times \cdots \times n_{d}$ rank-1 tensors $\mathfrak{a}=\left(\mathcal{A}_{1}, \ldots, \mathscr{A}_{r}\right)$ that satisfy:
(1) $\Phi_{r}(\mathfrak{a})$ is a smooth point of the semi-algebraic set $\sigma_{r}$;
(2) $\Phi_{r}(\mathfrak{a})$ is complex $r$-identifiable;
(3) the derivative $\mathrm{d}_{\mathfrak{a}} \Phi_{r}$ is injective

## Definition

The set of $r$-nice tensors is

$$
\mathcal{N}_{r ; \mathbf{n}}:=\Phi_{r}\left(\mathcal{M}_{r ; \mathbf{n}}\right)
$$

One can prove the following results:

## Proposition

Let $\sigma_{r}^{\mathbb{C}}$ be generically $r$-identifiable. Then, the following holds.
(1) $\mathcal{N}_{r, \mathrm{n}}$ is an open dense submanifold of $\sigma_{r}$.
(2) $\widehat{\mathcal{M}}_{r ; \mathbf{n}}:=\mathcal{M}_{r ; \mathbf{n}} / \mathfrak{S}_{r}$ is a manifold and the projection is a local diffeomorphism.
(3) The addition map

$$
\Phi_{r}: \widehat{\mathcal{M}}_{r ; \mathbf{n}} \rightarrow \mathcal{N}_{r ; \mathbf{n}},\left\{\mathcal{A}_{1}, \ldots, \mathscr{A}_{r}\right\} \rightarrow \mathcal{A}_{1}+\cdots+\mathcal{A}_{r}
$$

is a diffeomorphism.

## The condition number of the CPD

The inverse of $\Phi_{r}: \widehat{\mathcal{M}}_{r ; \mathbf{n}} \rightarrow \mathcal{N}_{r ; \mathbf{n}}$ is

$$
\tau_{r ; \mathbf{n}}: \mathcal{N}_{r ; \mathbf{n}} \rightarrow \widehat{\mathcal{M}}_{r ; \mathbf{n}}, \mathcal{A}_{1}+\cdots+\mathcal{A}_{r} \rightarrow\left\{\mathcal{A}_{1}, \ldots, \mathcal{A}_{r}\right\}
$$

We call it the tensor rank decomposition map.

## Definition

The condition number of the tensor rank decomposition for a tensor $\mathcal{A} \in \mathcal{N}_{r ; \mathbf{n}}$ is

$$
\kappa\left[\tau_{r ; \mathbf{n}}\right](\mathcal{A})=\left\|\mathrm{d}_{\mathcal{A}} \tau_{r ; \mathbf{n}}\right\|_{2}=\left\|\left(\mathrm{d}_{\mathfrak{a}} \Phi_{r}\right)^{-1}\right\|_{2}
$$

## Interpretation

If

$$
\begin{aligned}
& \mathcal{A}=\mathcal{A}_{1}+\cdots+\mathcal{A}_{r}=\sum_{i=1}^{r} \mathbf{a}_{i}^{1} \otimes \cdots \otimes \mathbf{a}_{i}^{d} \\
& \mathcal{B}=\mathcal{B}_{1}+\cdots+\mathcal{B}_{r}=\sum_{i=1}^{r} \mathbf{b}_{i}^{1} \otimes \cdots \otimes \mathbf{b}_{i}^{d}
\end{aligned}
$$

are tensors in $\mathbb{R}^{n_{1} \times \cdots \times n_{d}}$, then for $\|\mathcal{A}-\mathcal{B}\|_{F} \approx 0$ we have the asymptotically sharp bound

$$
\underbrace{\min _{\pi \in \mathfrak{S}_{r}} \sqrt{\sum_{i=1}^{r}\left\|\mathcal{A}_{i}-\mathcal{B}_{\pi_{i}}\right\|_{F}^{2}}}_{\text {forward error }} \lesssim \underbrace{\kappa\left[\tau_{r ; \mathbf{n}}\right](\mathcal{A})}_{\text {condition number }} \cdot \underbrace{\|\mathcal{A}-\mathcal{B}\|_{F}}_{\text {backward error }}
$$

## Back to our example

>> Ut\{1\} = randn $(25,25)$;
>> Ut\{2\} = randn $(25,25)$;
>> Ut\{3\} = randn $(25,25)$;
>> A = cpdgen(Ut);
>> Ur = cpd_gevd(A, 25);
>> E = kr(Ut) - kr(Ur);
>> norm( E(:), 2 ) / eps
ans =

$$
8.6249 \mathrm{e}+04
$$

We understand now that this can happen, because of a high condition number. However,
>> kappa = condition_number ( Ut )
ans =
2.134

The only explanation is that there is something wrong with the algorithm.

We show that algorithms based on a reduction to tensors in $\mathbb{R}^{n_{1} \times n_{2} \times 2}$ are numerically unstable.

The forward error produced by the algorithm divided by the backward error is "much" larger than the condition number, for some inputs.

## Pencil-based algorithms

A pencil-based algorithm (PBA) is an algorithm that computes the CPD of

$$
\mathcal{A}=\sum_{i=1}^{r} \mathbf{a}_{i} \otimes \mathbf{b}_{i} \otimes \mathbf{c}_{i} \in \mathcal{N}_{r ; n_{1}, n_{2}, n_{3}} \subset \mathbb{R}^{n_{1} \times n_{2} \times n_{d}}
$$

in the following way:
S1. Choose a fixed $Q \in \mathbb{R}^{n_{3} \times 2}$ with orthonormal columns.
S2. $\mathcal{B} \leftarrow\left(I, I, Q^{T}\right) \cdot \mathfrak{A}$;
S3. $\left\{\mathbf{a}_{1}, \ldots, \mathbf{a}_{r}\right\} \leftarrow \widehat{\theta}(\mathcal{B})$;
S4.a Choose an order $A:=\left(\mathbf{a}_{1}, \ldots, \mathbf{a}_{r}\right)$;
S4.b $\left(\mathbf{b}_{1} \otimes \mathbf{c}_{1}, \ldots, \mathbf{b}_{r} \otimes \mathbf{c}_{r}\right) \leftarrow\left(A^{\dagger} \mathcal{A}_{(1)}\right)^{T}$;
S5. output $\leftarrow \pi\left(\odot\left(\left(\mathbf{a}_{1}, \ldots, \mathbf{a}_{r}\right),\left(\mathbf{b}_{1} \otimes \mathbf{c}_{1}, \ldots, \mathbf{b}_{r} \otimes \mathbf{c}_{r}\right)\right)\right)$.
where $\pi: \mathcal{S}^{\times r} \rightarrow\left(\mathcal{S}^{\times r} / \mathfrak{S}_{r}\right)$ and $\odot$ is the Khatri-Rao product:
$\odot(A, B):=\left(\mathbf{a}_{i} \otimes \mathbf{b}_{i}\right)_{i}$.

## Pencil-based algorithms

A pencil-based algorithm (PBA) is an algorithm that computes the CPD of

$$
\mathcal{A}=\sum_{i=1}^{r} \mathbf{a}_{i} \otimes \mathbf{b}_{i} \otimes \mathbf{c}_{i} \in \mathcal{N}_{r ; n_{1}, n_{2}, n_{3}} \subset \mathbb{R}^{n_{1} \times n_{2} \times n_{d}}
$$

in the following way:
OK Choose a fixed $Q \in \mathbb{R}^{n_{3} \times 2}$ with orthonormal columns.
OK $\quad \mathcal{B} \leftarrow\left(I, I, Q^{T}\right) \cdot \mathfrak{A}$;
BAD $\left\{\mathbf{a}_{1}, \ldots, \mathbf{a}_{r}\right\} \leftarrow \widehat{\theta}(\mathcal{B})$;
OK Choose an order $A:=\left(\mathbf{a}_{1}, \ldots, \mathbf{a}_{r}\right)$;
OK $\quad\left(\mathbf{b}_{1} \otimes \mathbf{c}_{1}, \ldots, \mathbf{b}_{r} \otimes \mathbf{c}_{r}\right) \leftarrow\left(A^{\dagger} \mathcal{A}_{(1)}\right)^{T}$;
OK output $\leftarrow \pi\left(\odot\left(\left(\mathbf{a}_{1}, \ldots, \mathbf{a}_{r}\right),\left(\mathbf{b}_{1} \otimes \mathbf{c}_{1}, \ldots, \mathbf{b}_{r} \otimes \mathbf{c}_{r}\right)\right)\right)$.
where $\pi: \mathcal{S}^{\times r} \rightarrow\left(\mathcal{S}^{\times r} / \mathfrak{S}_{r}\right)$ and $\odot$ is the Khatri-Rao product:
$\odot(A, B):=\left(\mathbf{a}_{i} \otimes \mathbf{b}_{i}\right)_{i}$.

The magic map $\hat{\theta}$ needs to recover the vectors from the first factor matrix when restricted to $\mathcal{N}_{r ; n_{1}, n_{2}, 2}$ :

$$
\begin{aligned}
\left.\widehat{\theta}\right|_{\mathcal{N}_{r ; n_{1}, n_{2}, 2}}: & \mathcal{N}_{r ; n_{1}, n_{2}, 2}
\end{aligned} \longrightarrow^{\longrightarrow}\left(\mathbb{S}^{n_{1}-1}\right)^{\times r} / \mathfrak{S}_{r} .
$$

Since the input to $\widehat{\theta}$ will be the result of a previous numerical computation, the domain of definition of $\hat{\theta}$ should also encompass a sufficiently large neighborhood of $\mathcal{N}_{r ; n_{1}, n_{2}, 2}$ !

For proving instability, it does not matter what $\widehat{\theta}$ computes outside of $\mathcal{N}_{r ; n_{1}, n_{2}, 2}$.

For a valid input $\mathcal{A} \in \mathcal{N}_{r ; n_{1}, n_{2}, n_{3}}$, let $\left\{\tilde{\mathscr{A}}_{1}, \ldots, \tilde{\mathscr{A}}_{r}\right\}$ be the CPD (in floating-point representation) returned by the PBA.

Our proof strategy consists of showing that for every $\epsilon>0$ there exists an open neighborhood $\mathcal{O}_{\epsilon} \subset \mathcal{N}_{r ; n_{1}, n_{2}, n_{3}}$ of $r$-nice tensors such that the excess factor

$$
\begin{aligned}
\omega(\mathcal{A}) & =\frac{\text { observed forward error due to algorithm }}{\text { maximum forward error due to problem }} \\
& :=\frac{\min _{\pi \in \mathfrak{S}_{r}} \sqrt{\sum_{i=1}^{r}\left\|\mathcal{A}_{i}-\widetilde{\mathcal{A}}_{i}\right\|^{2}}}{\kappa\left[\tau_{r ; n_{1}, n_{2}, n_{3}}\right](\mathcal{A}) \cdot\|\mathcal{A}-\mathrm{f}(\mathscr{A})\|_{F}}
\end{aligned}
$$

behaves like a constant times $\epsilon^{-1}$.

Formally, we showed the following result:

## Theorem (Beltrán, Breiding, Vannieuwenhoven (2018))

There exist a constant $k>0$ and a tensor $O \in \mathcal{N}_{r ; n_{1}, n_{2}, n_{3}}$ with the following properties: For all sufficiently small $\epsilon>0$, there exists an open neighborhood $\mathcal{O}_{\epsilon}$ of $O$, such that for all tensors $\mathfrak{A} \in \mathcal{O}_{\epsilon}$ we have
(1) $\mathcal{A} \in \mathcal{N}_{r ; n_{1}, n_{2}, n_{3}}$ is a valid input for a PBA, and
(2) $\omega(\mathscr{A}) \geq k \epsilon^{-1}$.

In other words, the forward error produced by a PBA can be larger than the maximum forward error expected from the tensor decomposition problem by an arbitrarily large factor.

## Distribution of the condition number

We conceived the existence of this problem after seeing the distribution of the condition number of random rank-1 tensors

$$
\mathcal{A}_{i}=\mathbf{a}_{i} \otimes \mathbf{b}_{i} \otimes \mathbf{c}_{i} \in \mathbb{R}^{n_{1} \times n_{2} \times n_{3}}
$$

where

- $\mathbf{a}_{i} \in \mathbb{R}^{n_{1}}$ and $\mathbf{b}_{i} \in \mathbb{R}^{n_{2}}$ are arbitrary, and
- the $\mathbf{c}_{i} \in \mathbb{R}^{n_{3}}$ are i.i.d. random Gaussian vectors.

We showed, based on Cai, Fan, and Jiang (2013), that

$$
\mathrm{P}[\kappa \geq \alpha] \geq \mathrm{P}\left[\max _{1 \leq i \neq j \leq r} \frac{1}{\sqrt{1-\left\langle\mathbf{c}_{i}, \mathbf{c}_{j}\right\rangle}} \geq \alpha\right] \rightarrow 1-e^{-K r^{2} \alpha^{1-n_{3}}}
$$

as $r \rightarrow \infty$; herein, $K$ is a constant depending only on $n_{3}$.

## Empirical distribution of $\kappa(\mathcal{A})$


$10^{5}$ random rank-15 tensors $\sum_{i=1}^{15} \mathbf{a}_{i} \otimes \mathbf{b}_{i} \otimes \mathbf{c}_{i} \in \mathbb{R}^{15 \times 15 \times n}$, where

$$
A=\left[\mathbf{a}_{1}, \ldots, \mathbf{a}_{15}\right] \text { and } B=\left[\mathbf{b}_{1}, \ldots, \mathbf{b}_{15}\right] \text { and } C=\left[\mathbf{c}_{1}, \ldots, \mathbf{c}_{15}\right]
$$

are Gaussian matrices.
dashed lines $=$ empirical distribution.

## Empirical distribution of $\kappa(\mathcal{A})$


$10^{5}$ random rank-15 tensors $\sum_{i=1}^{15} \mathbf{a}_{i} \otimes \mathbf{b}_{i} \otimes \mathbf{c}_{i} \in \mathbb{R}^{15 \times 15 \times n}$, where

$$
A=\left[\mathbf{a}_{1}, \ldots, \mathbf{a}_{15}\right] \text { and } B=\left[\mathbf{b}_{1}, \ldots, \mathbf{b}_{15}\right]
$$

are random orthogonal matrices and $C=\left[\mathbf{c}_{1}, \ldots, \mathbf{c}_{15}\right]$ is a Gaussian matrix.
dashed lines = empirical distribution.

## Empirical distribution of the excess factor


$10^{5}$ random rank-17 tensors $\sum_{i=1}^{17} \mathbf{a}_{i} \otimes \mathbf{b}_{i} \otimes \mathbf{c}_{i} \in \mathbb{R}^{23 \times 17 \times 5}$ :

$$
A=\left[\mathbf{a}_{1}, \ldots, \mathbf{a}_{17}\right] \text { and } B=\left[\mathbf{b}_{1}, \ldots, \mathbf{b}_{17}\right] \text { and } C=\left[\mathbf{c}_{1}, \ldots, \mathbf{c}_{17}\right]
$$

are Gaussian matrices.

## Empirical distribution of the excess factor


$10^{5}$ random rank-17 tensors $\sum_{i=1}^{17} \mathbf{a}_{i} \otimes \mathbf{b}_{i} \otimes \mathbf{c}_{i} \in \mathbb{R}^{23 \times 17 \times 15}$ :

$$
A=\left[\mathbf{a}_{1}, \ldots, \mathbf{a}_{17}\right] \text { and } B=\left[\mathbf{b}_{1}, \ldots, \mathbf{b}_{17}\right] \text { and } C=\left[\mathbf{c}_{1}, \ldots, \mathbf{c}_{17}\right]
$$

are Gaussian matrices.

## Empirical distribution of the excess factor


$10^{5}$ random rank- 23 tensors $\sum_{i=1}^{23} \mathbf{a}_{i} \otimes \mathbf{b}_{i} \otimes \mathbf{c}_{i} \in \mathbb{R}^{23 \times 23 \times 23}$ :

$$
A=\left[\mathbf{a}_{1}, \ldots, \mathbf{a}_{23}\right] \text { and } B=\left[\mathbf{b}_{1}, \ldots, \mathbf{b}_{23}\right] \text { and } C=\left[\mathbf{c}_{1}, \ldots, \mathbf{c}_{23}\right]
$$

are Gaussian matrices.

## Conclusions

Take-away story:
(1) Tensors are conjectured to be generically $r$-identifiable for almost all low ranks $r$.
(2) The condition number of the CPD measures the stability of the unique rank-1 tensors.
(3) Reduction to a matrix pencil yields numerically unstable algorithms for computing CPDs.

## Further reading

- Beltrán, Breiding, and Vannieuwenhoven, Pencil-based algorithms for tensor rank decomposition are not stable, SIAM J. Matrix Anal. and Appl.
- Beltrán, Breiding, and Vannieuwenhoven, The average condition number of most tensor rank decomposition problems is infinite, arXiv1903.05527.
- Breiding and Vannieuwenhoven, The condition number of join decompositions, SIAM J. Matrix Anal. and Appl., 2018.
- Breiding and Vannieuwenhoven, On the average condition number of tensor rank decompositions, arXiv:1801.01673.
- Breiding and Vannieuwenhoven, A Riemannian trust region method for the canonical tensor rank approximation problem, SIAM J. Optim, 2018.

- Kruskal, Three-way arrays: rank and uniqueness of trilinear decompositions, with application to arithmetic complexity and statistics, Lin. Alg. Appl., 1977.
- Abo, Ottaviani and Peterson, Induction for secant varieties of Segre varieties, Trans. AMS, 2009.
- Chiantini and Ottaviani, On Generic Identifiability of 3-Tensors of Small Rank, SIMAX, 2012.
- Bocci and Chiantini, On the identifiability of binary Segre products, J. Algebraic Geom., 2013.
- Bocci, Chiantini, and Ottaviani, Refined methods for the identifiability of tensors, Ann. Mat. Pura Appl., 2013.
- Chiantini, Mella and Ottaviani, One example of general unidentifiable tensors, J. Alg. Stat., 2014.
- Chiantini, Ottaviani, and Vannieuwenhoven, An algorithm for generic and low-rank specific identifiability of complex tensors, SIMAX, 2014.
- Qi, Comon, and Lim, Semialgebraic geometry of nonnegative tensor rank, SIMAX, 2016.

